



360搜索技术论坛长期交流群





分布式存储实践

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O. 介绍

- 2013年加入360
- 主要项目
 - Bada 类Dynamo 分布式存储
 - Pika 兼容Redis 协议Nosql
 - Zeppelin 海量数据存储, 支持pb, s3协议
 - 基础库: Pink, Slash, floyd 等等



O. 大纲

- 在线存储
- 了解硬件
- 基础理论
- 存储引擎
- zeppelin 设计



O. 离线存储

- 离线存储
 - 吞吐
 - 成本



O. 在线存储

- 在线存储
 - 性能
 - 稳定性
 - 接口



O. 在线存储

- key-value
- 表格存储
- 块存储
- S3(Simple Storage Service)
- 图数据库
- new-sql
- ...



O. Bada



- key-value store
- 云盘场景: 所有云盘用户文件索引, 通过bada查询文件所在存储节点 写入近百亿, 峰值QPS20万, 数据总量千亿, 单条value 64字节, 数据总量20T
- onebox场景: 存储关键字的rank 信息在bada中, 也无需要在4个机房数据同步, 单机房亿级别访问量

O. ceph

- 块存储
- 虚拟机提供块存储服务
- 线上500 台虚拟机



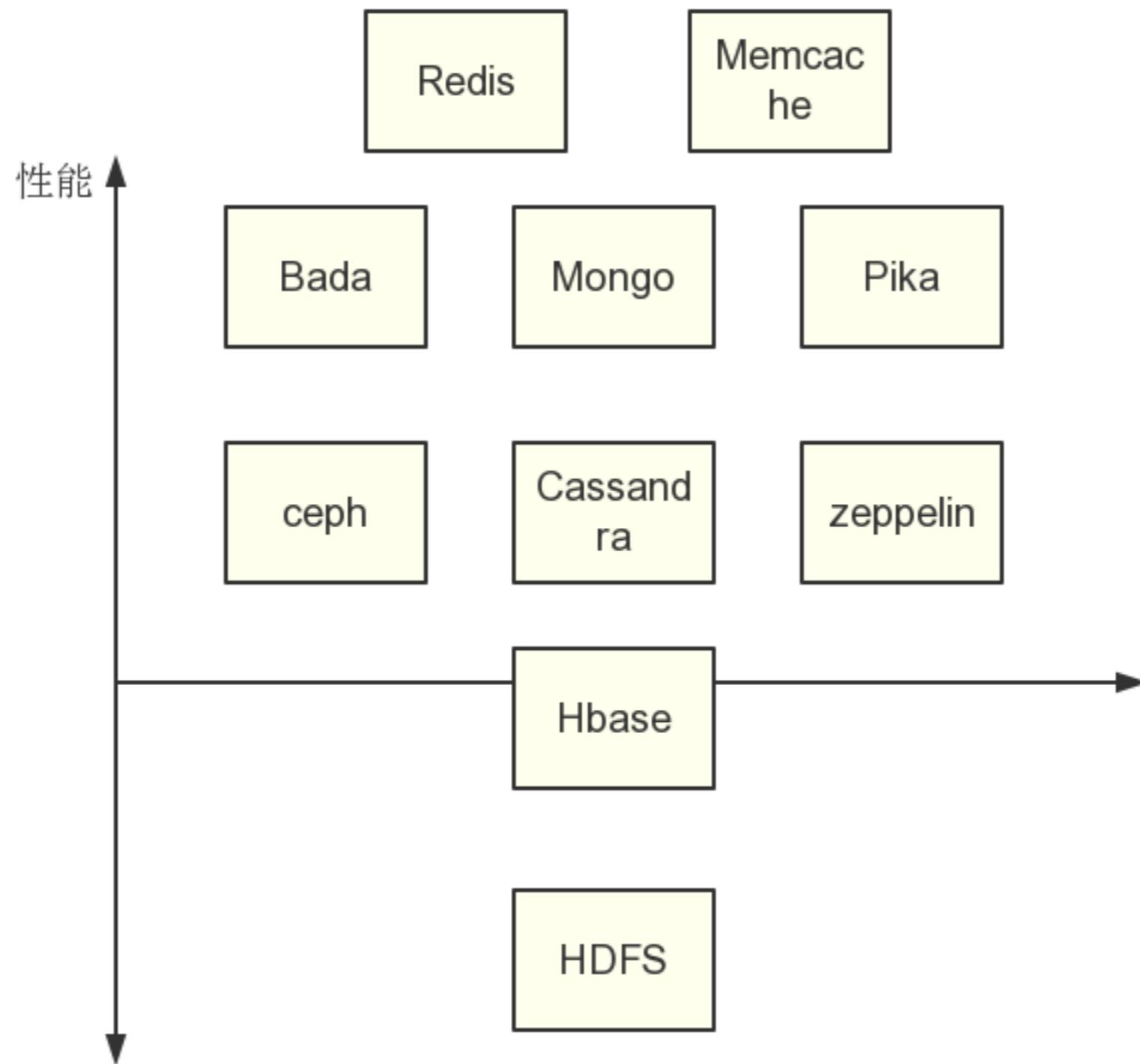
- 大容量, 支持redis 接口的Nosql
- 360 内部500+ 实例, 数据大小10T.
- 数据分析业务之前7个redis 集群, 每个20G, 每次分析挨个遍历redis 集群. 迁移pika 以后非常方便
- 手助存用户下载信息库, 历史遗留问题存在redis 30G. 担心主库挂重灌对用户影响. 迁移pika 以后故障秒切
- 外部公司大量使用

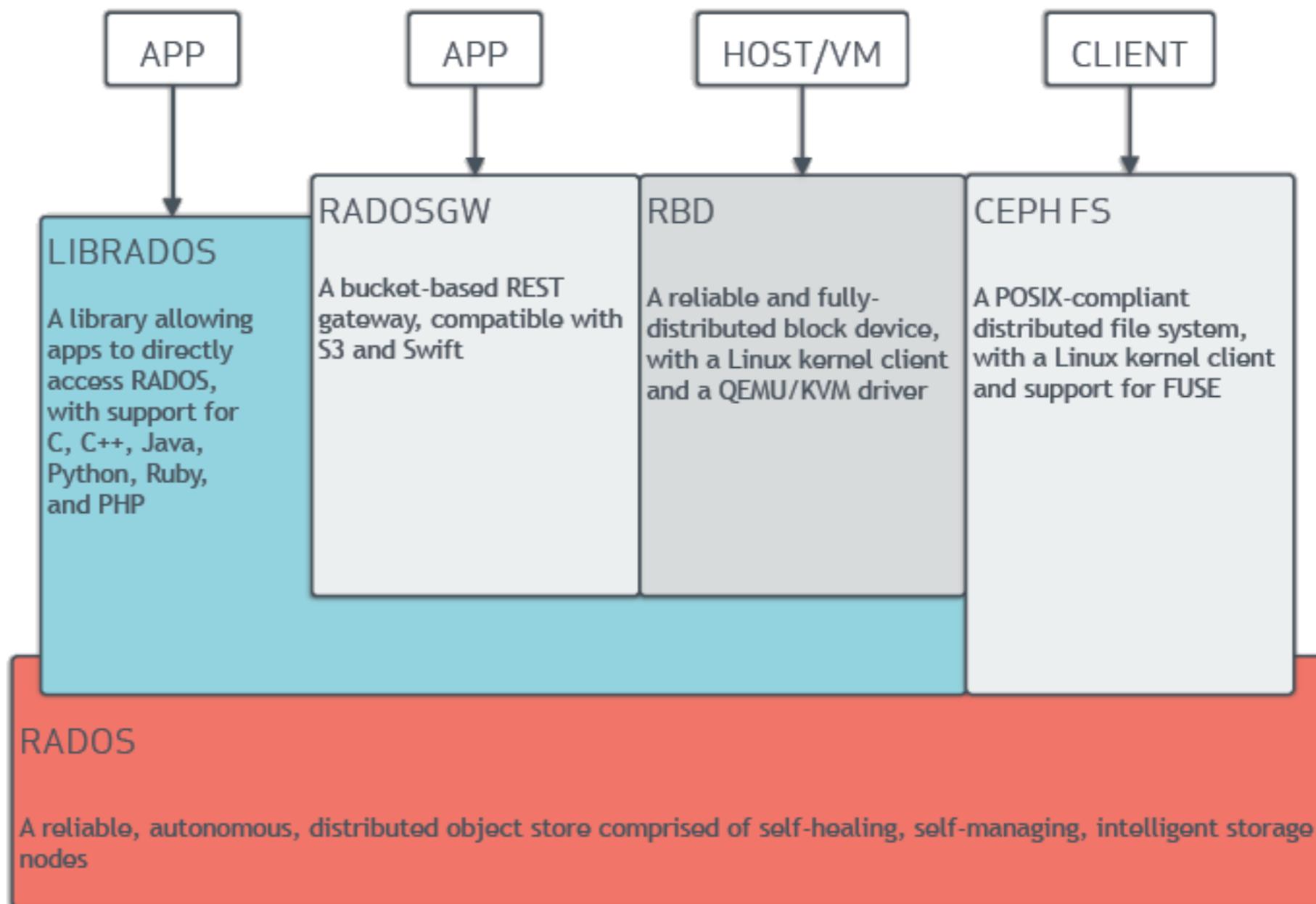


O. zeppelin



- key-value store, S3(Simple Storage Service)
- 存储PB 级别数据
- 场景: docker 容器镜像存储
- 低成本





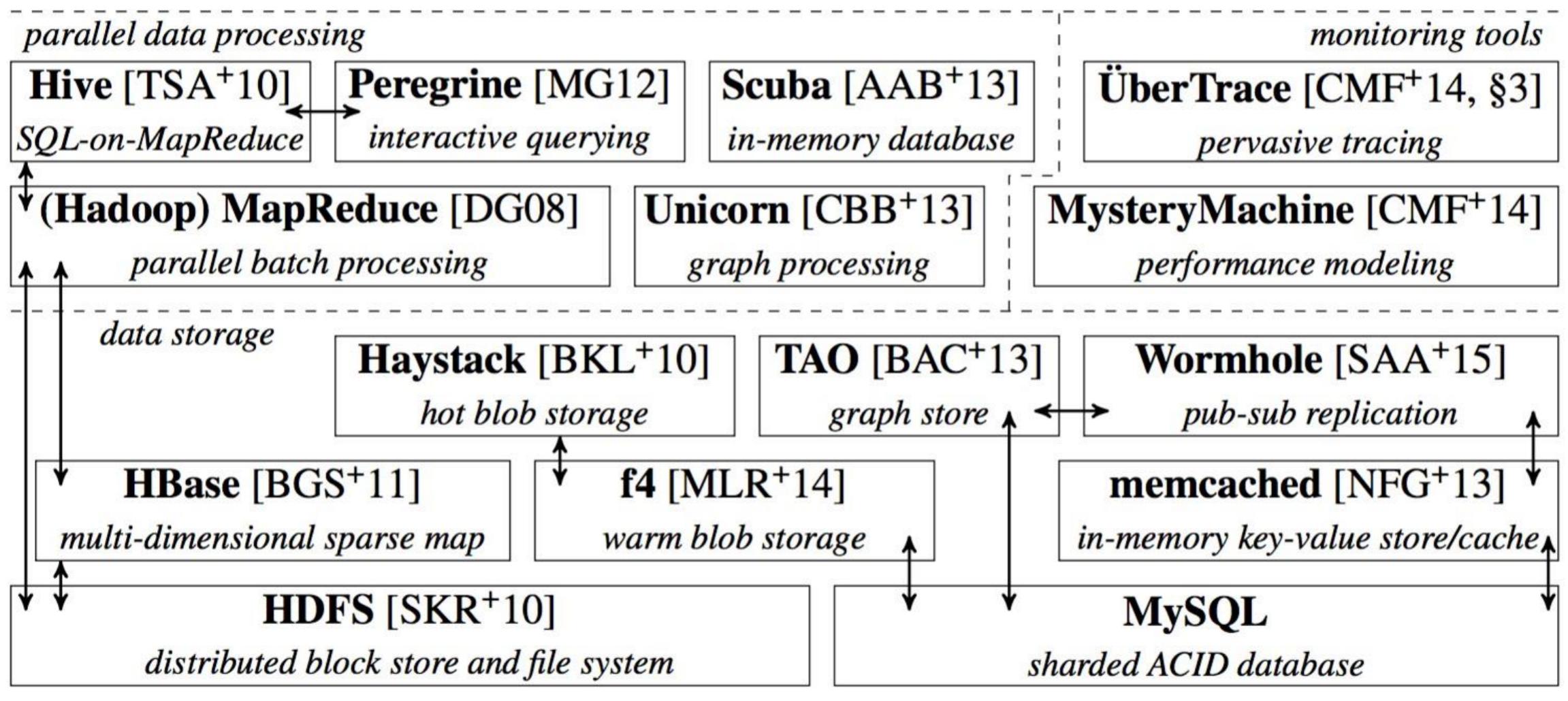


Figure 1: The Facebook infrastructure stack. I omit front-end serving systems about which details are unknown. Arrows indicate data exchange and dependencies between systems; simple layering does *not* imply a dependency or relation.

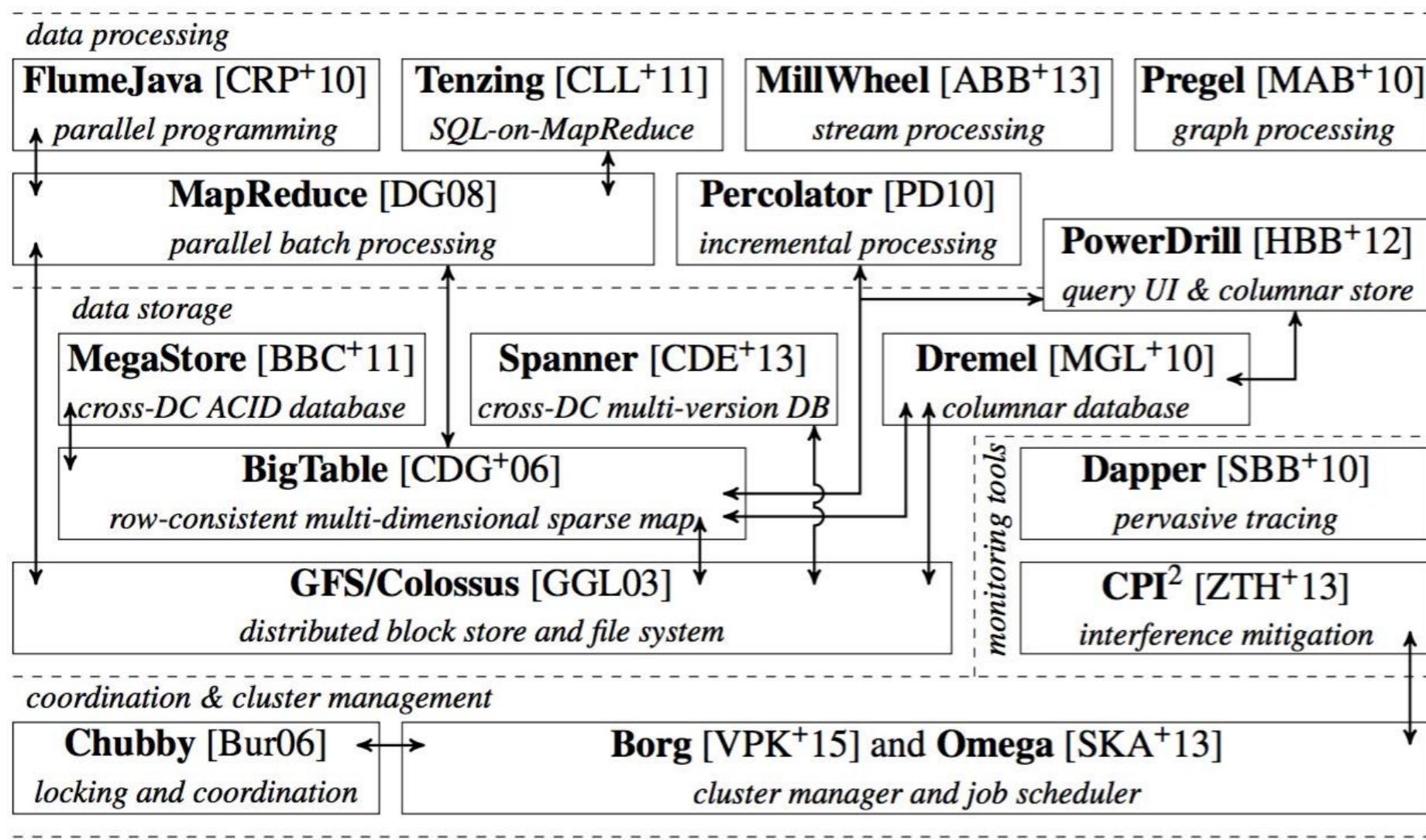


Figure 1: The Google infrastructure stack. I omit the F1 database [SOE⁺12] (the back-end of which was superseeded by Spanner), and unknown front-end serving systems. Arrows indicate data exchange and dependencies between systems; simple layering does *not* imply a dependency or relation.



Performance

Performance

1. Response time: 完成一个任务需要的时间
2. Throughput: 指定时间内可以完成的任务数
3. Average VS Quantile

1

**Two Lists of Response Times
(average for each is 1.000 second)**

	List A	List B
1	.924	.796
2	.928	.798
3	.954	.802
4	.957	.823
5	.961	.919
6	.965	.977
7	.972	1.076
8	.979	1.216
9	.987	1.273
10	1.373	1.320



Numbers Everyone Should Know

L1 cache reference	0.5 ns
Branch mispredict	5 ns
L2 cache reference	7 ns
Mutex lock/unlock	25 ns
Main memory reference	100 ns
Compress 1K bytes with Zippy	3,000 ns
Send 2K bytes over 1 Gbps network	20,000 ns
Read 1 MB sequentially from memory	250,000 ns
Round trip within same datacenter	500,000 ns
Disk seek	10,000,000 ns
Read 1 MB sequentially from disk	20,000,000 ns
Send packet CA->Netherlands->CA	150,000,000 ns

Jeff Dean [Designs, Lessons and Advice from Building Large Distributed Systems](#)



CATEGORY	REPRESENTATIVE DEVICE	SEQUENTIAL READ BANDWIDTH	SEQUENTIAL WRITE BANDWIDTH	4KB READ IOPS	4KB WRITE IOPS
Mechanical disk	Western Digital Black WD4001FAEX (4TB)	130MB/s	130MB/s	110	150
SATA-attached SSD	Samsung 850 Pro (1TB)	550MB/s	520MB/s	10,000	36,000
PCIe-attached SSD	Intel 750 (1.2TB)	2,400MB/s	1,200MB/s	440,000	290,000
Main memory	Skylake @ 3200MHz	42,000MB/s	48,000MB/s	16,100,000 (62ns/operation)	

(In the above table, all IOPS figures are reported assuming a queue depth of 1, so will tend to be worst case numbers for the SSDs.)



CATEGORY	IMPLIED SEEK TIME	IMPLIED SEEK TIME	MEAN IMPLIED
	FROM READ	FROM WRITE	SEEK TIME
Mechanical Disk	9.06ms	6.63ms	7.85ms
SATA-at- tached SSD	92.8us	20.2us	56.5us
PCIe-at- tached SSD	645ns	193ns	419ns

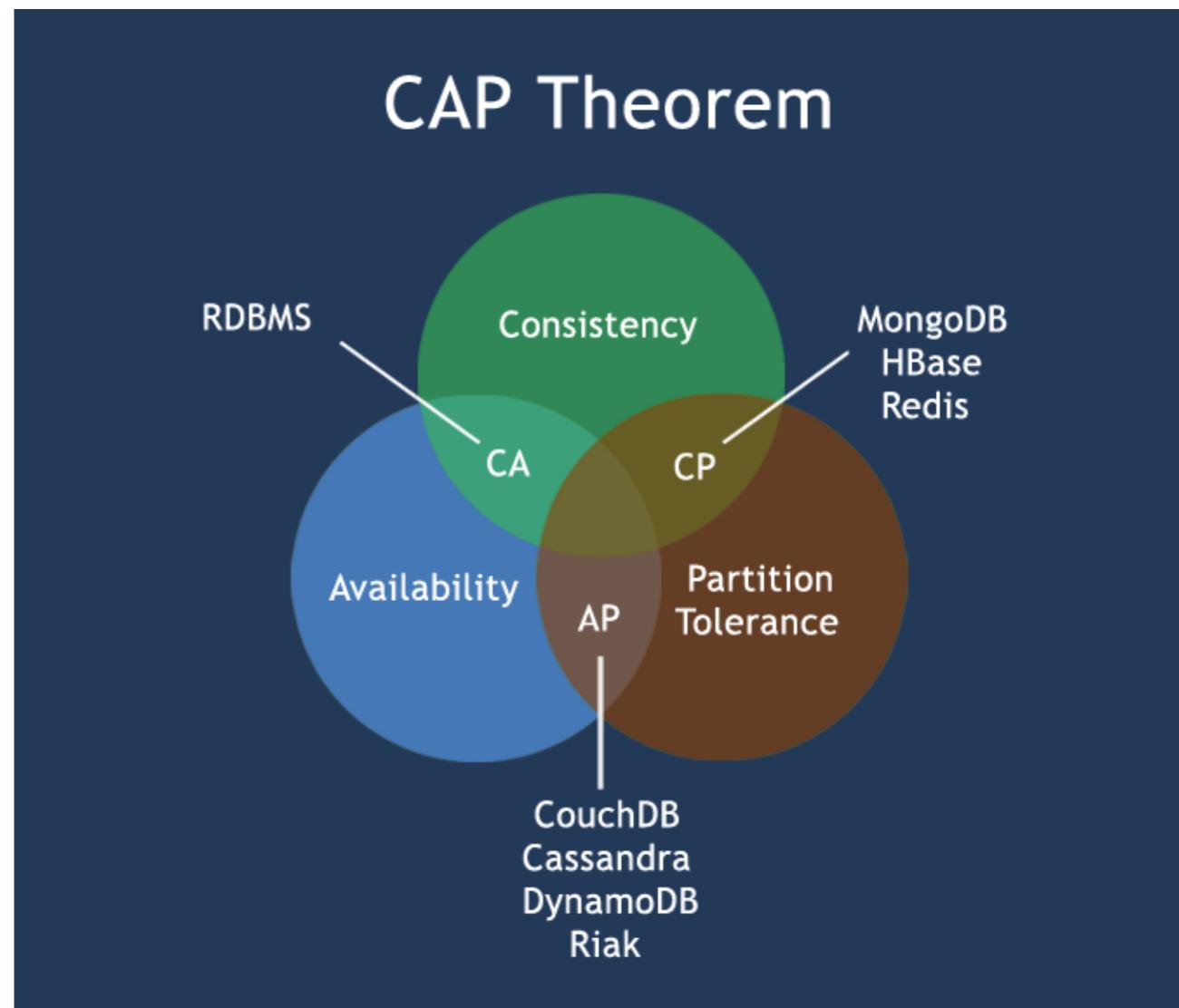
O. Example



- 对于SATA 盘的zeppelin, 写入一条256K 大小数据延迟
 - $9.06 \text{ ms} + 256\text{k} * 20 \text{ ms/MB} = 14.06 \text{ ms}$
- bada 北京机房与上海机房之间同步延迟极限
 - $1500\text{km} / (3 * 10^8) * 2 = 10\text{ms}$

CAP

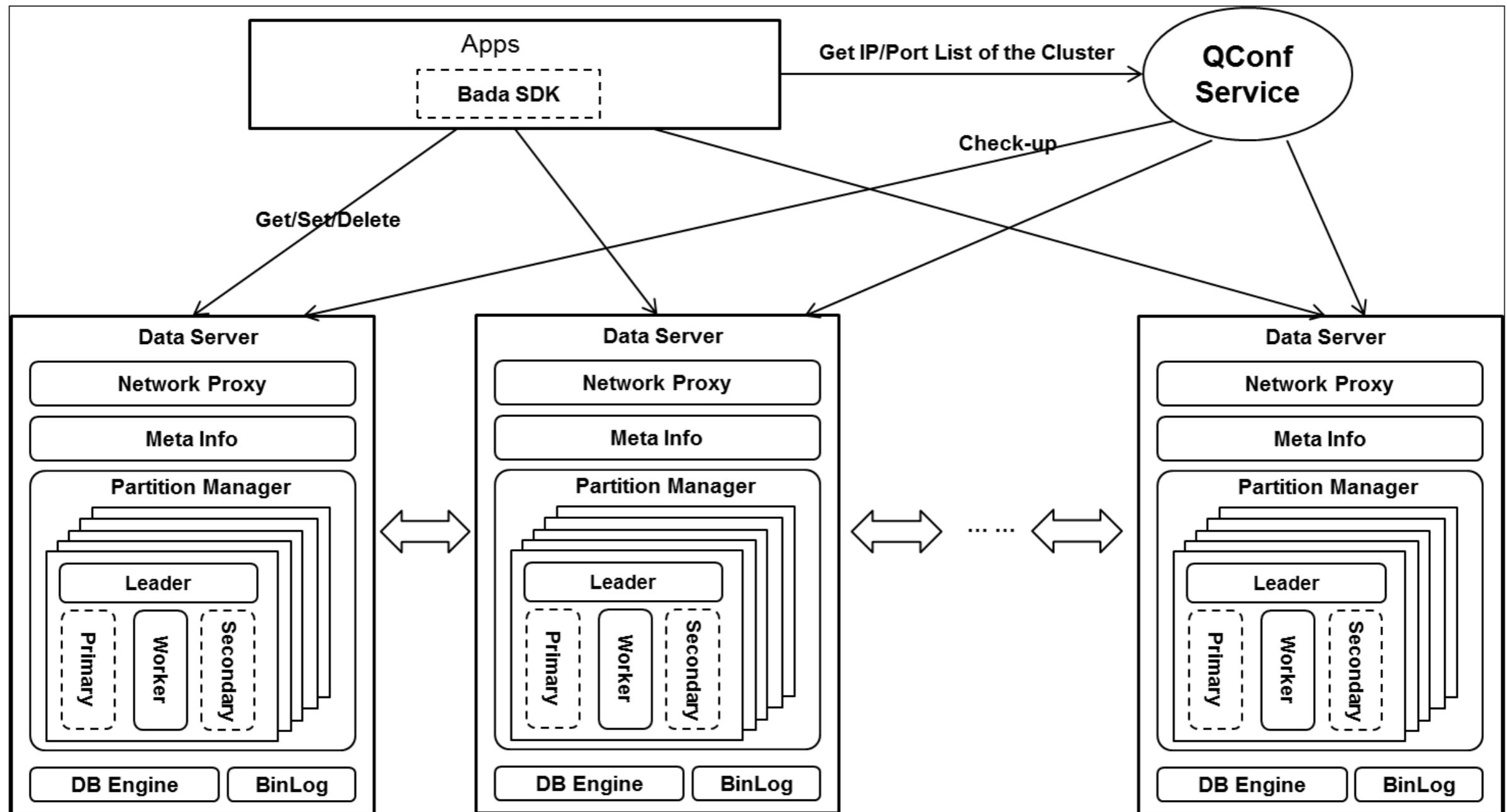
Eric Brewer 1998年提出
Consistency
Availability
Partition Tolerance

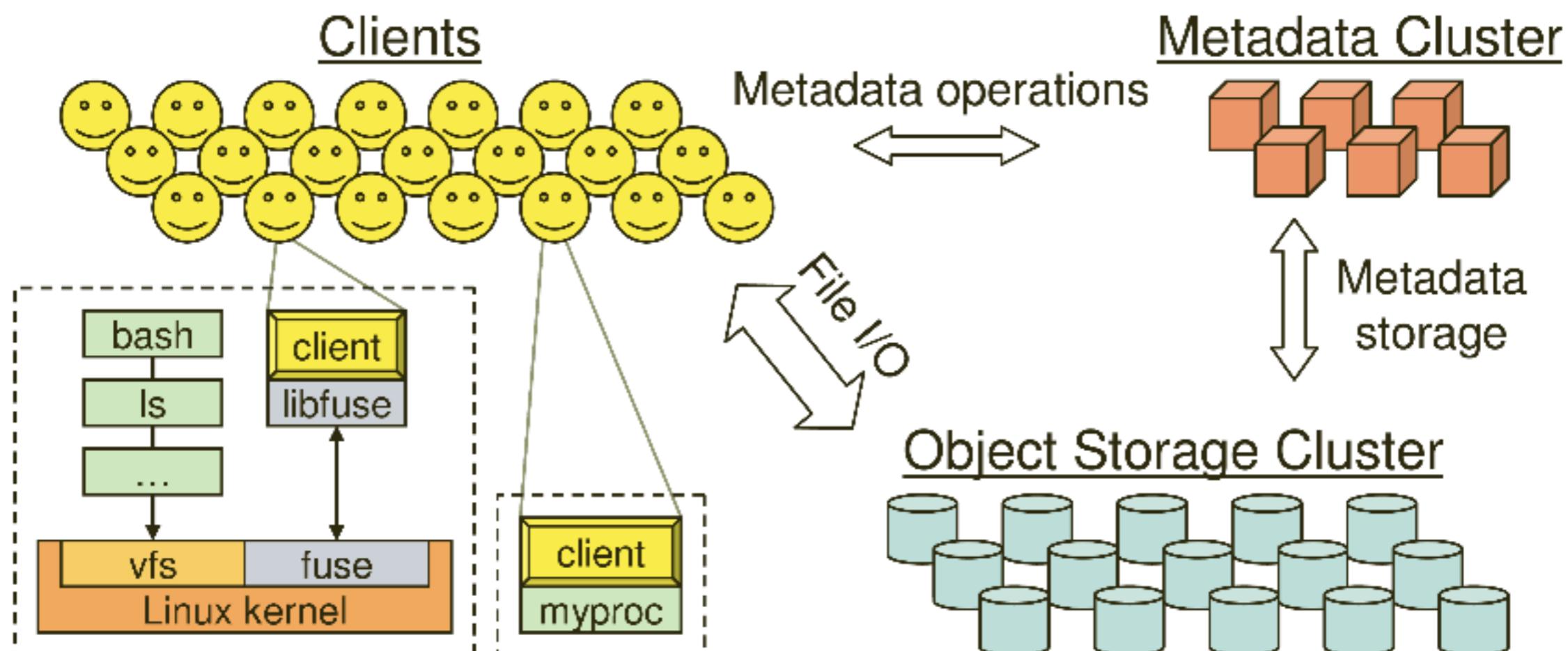


O. 基础理论

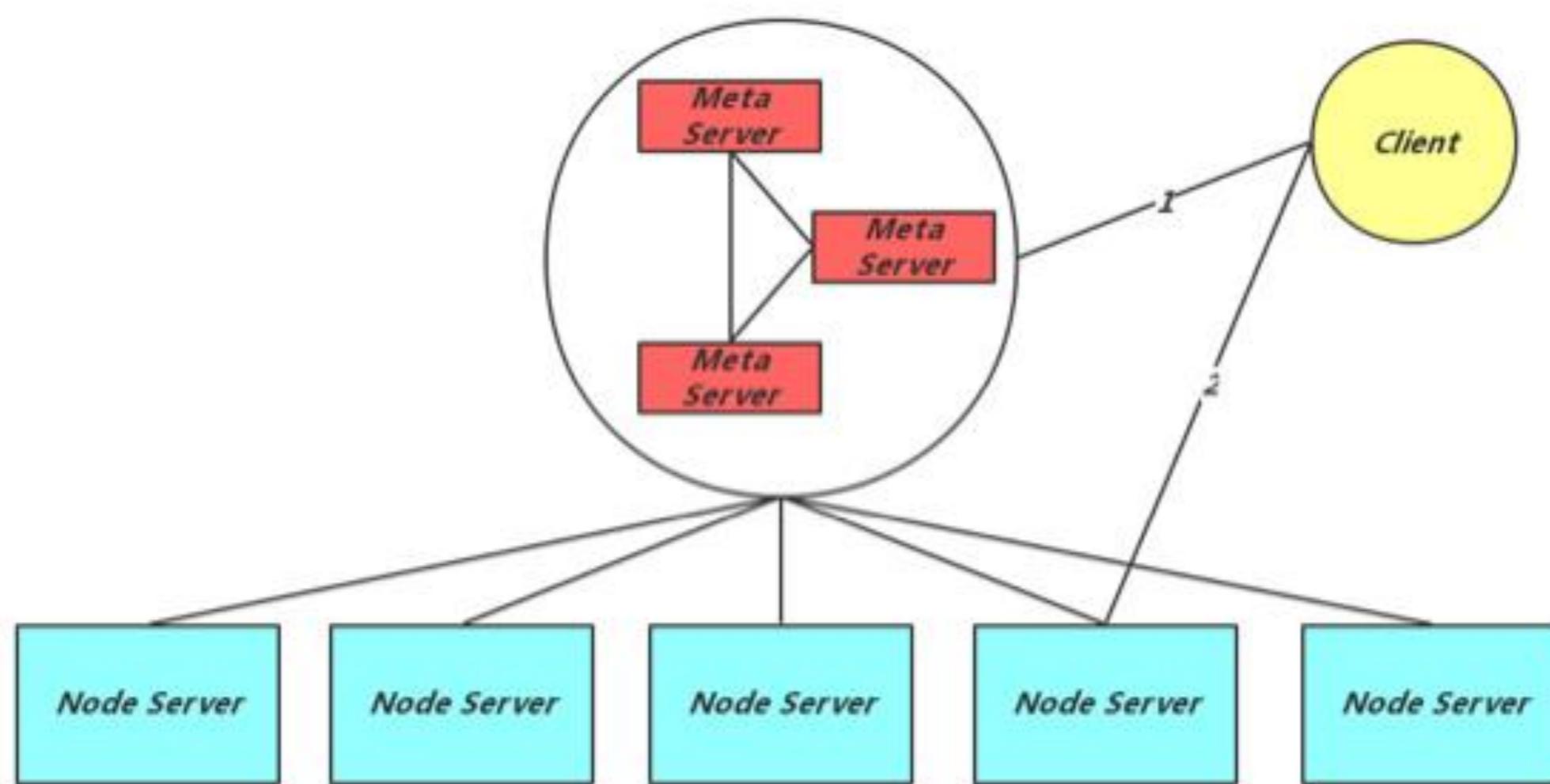


- 中心节点设计
 - 有中心 VS 去中心
- 副本策略
 - 3副本主从 VS Quorum VS erasure code
- 数据分布策略
 - consistent hash vs range map
- 一致性协议
 - 弱一致, 最终一致, 强一致
 - multi-paxos vs raft





O. zeppelin



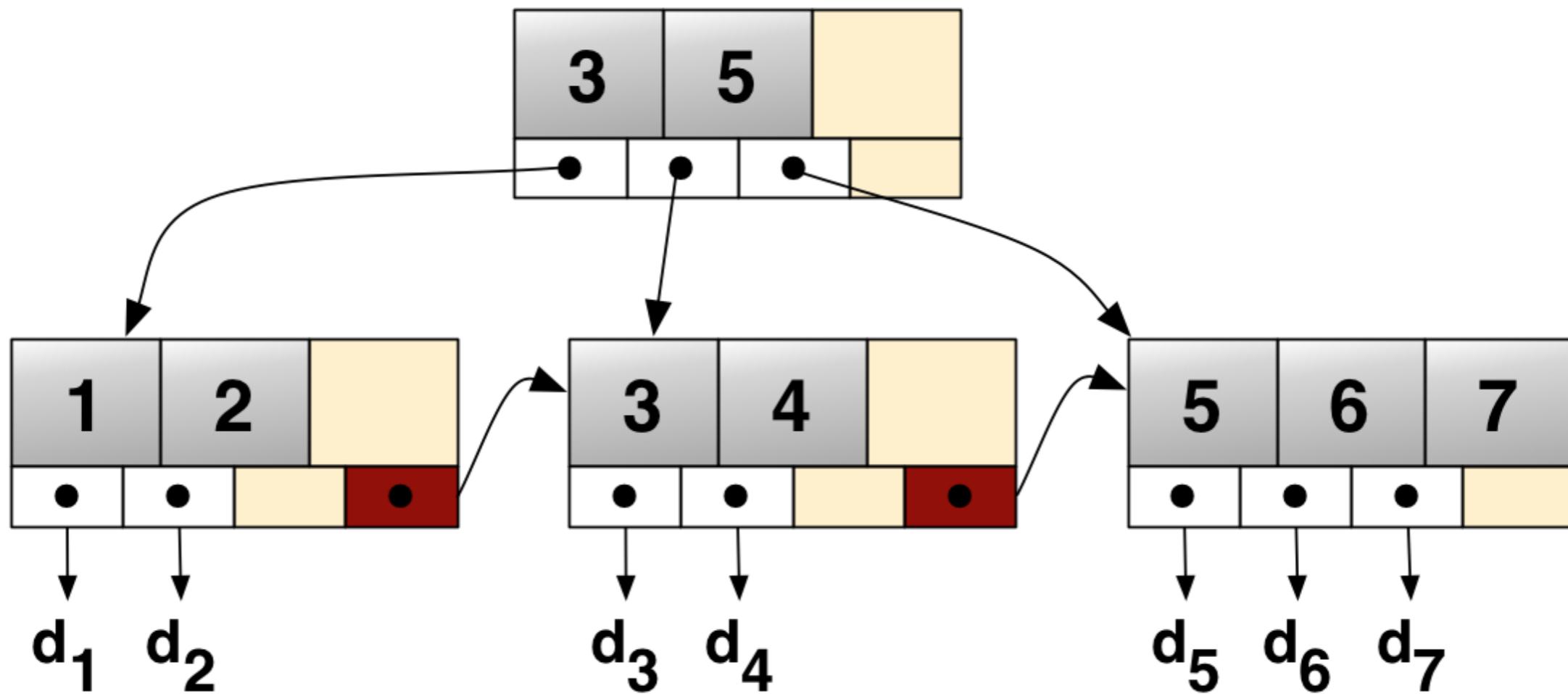


存储引擎

O. B+ Tree

360搜索
so.com

IT大咖说
不止于技术



O. LSM Tree

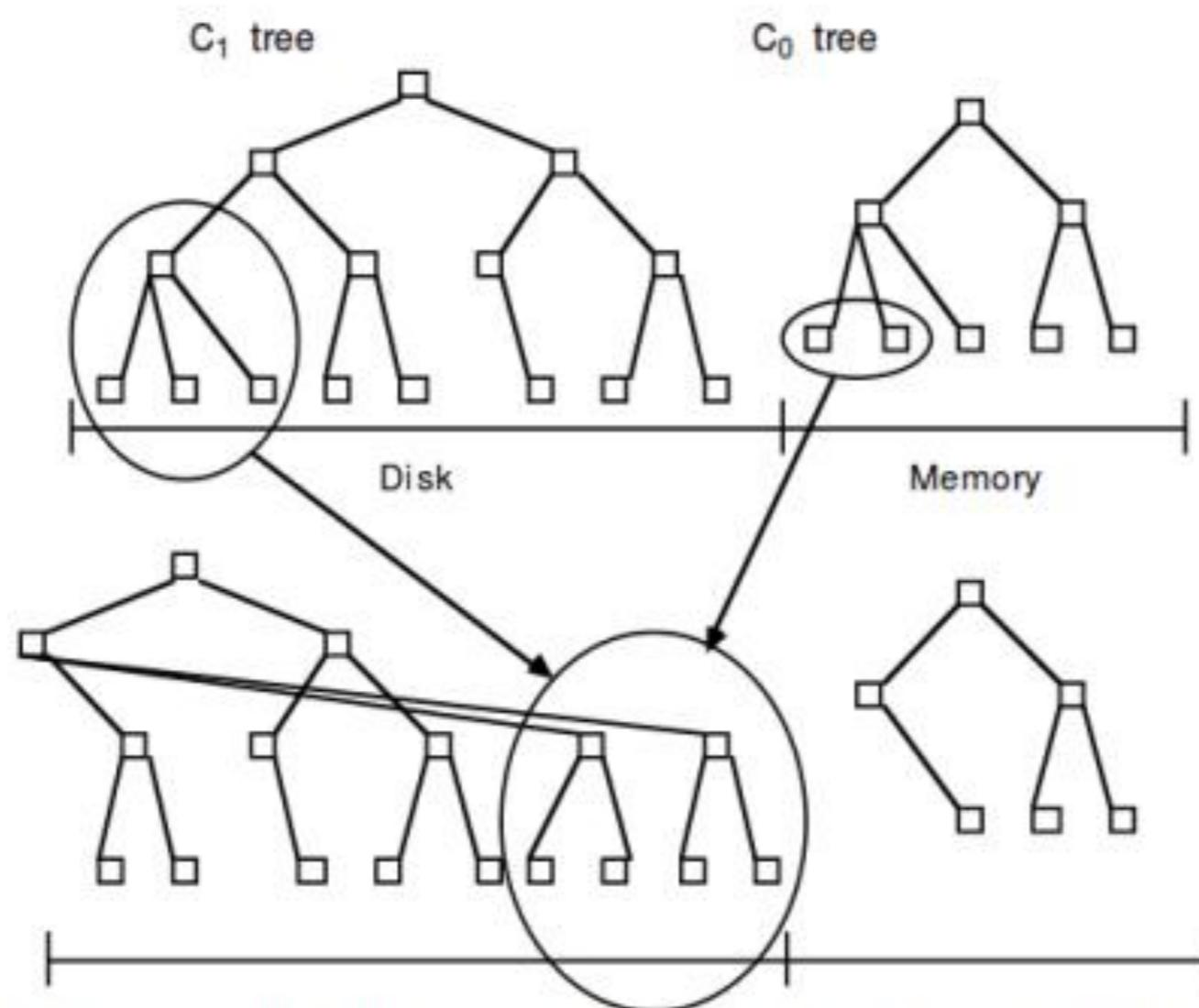
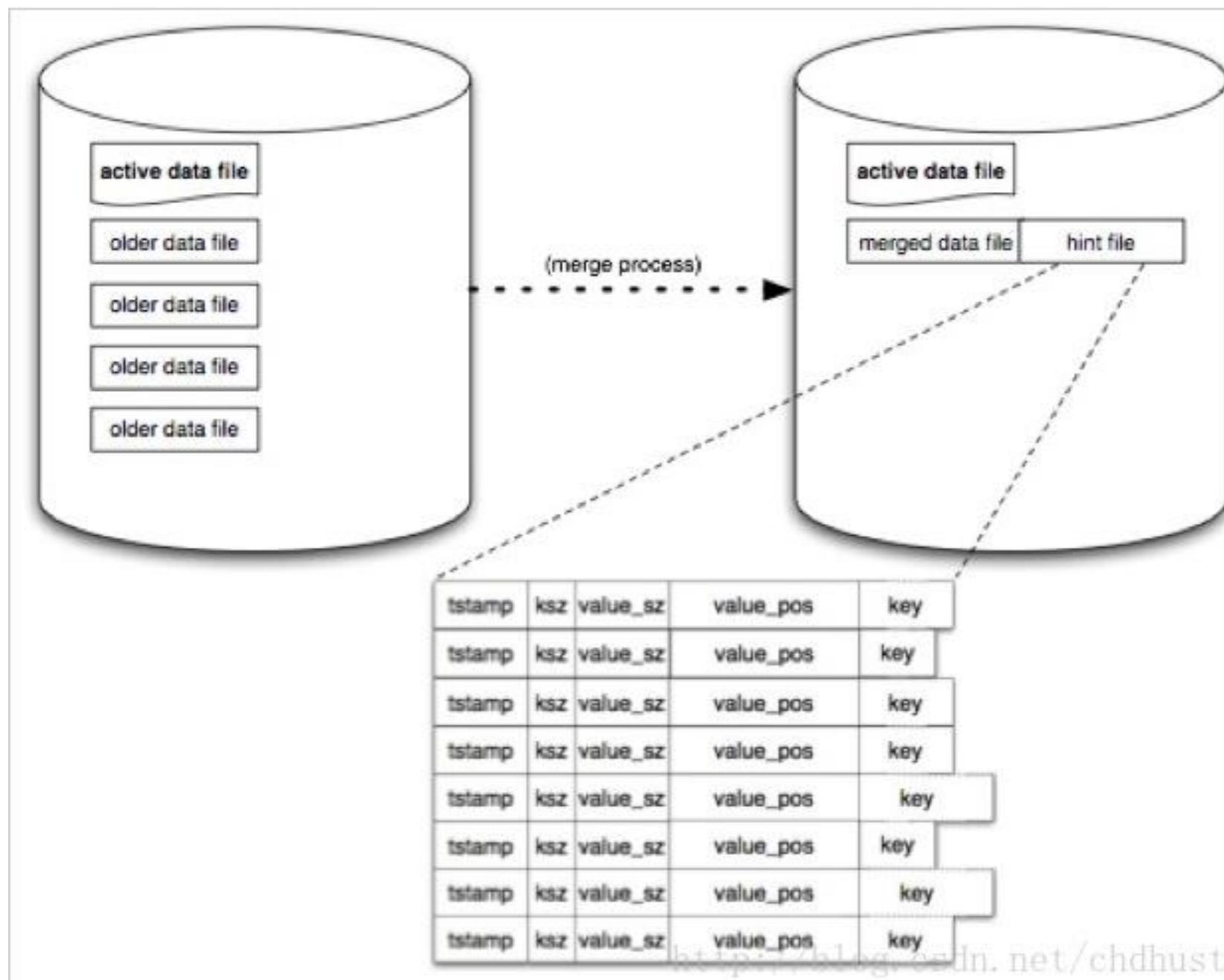


Figure 2.2. Conceptual picture of rolling merge steps, with result written back to disk

O. bitcask



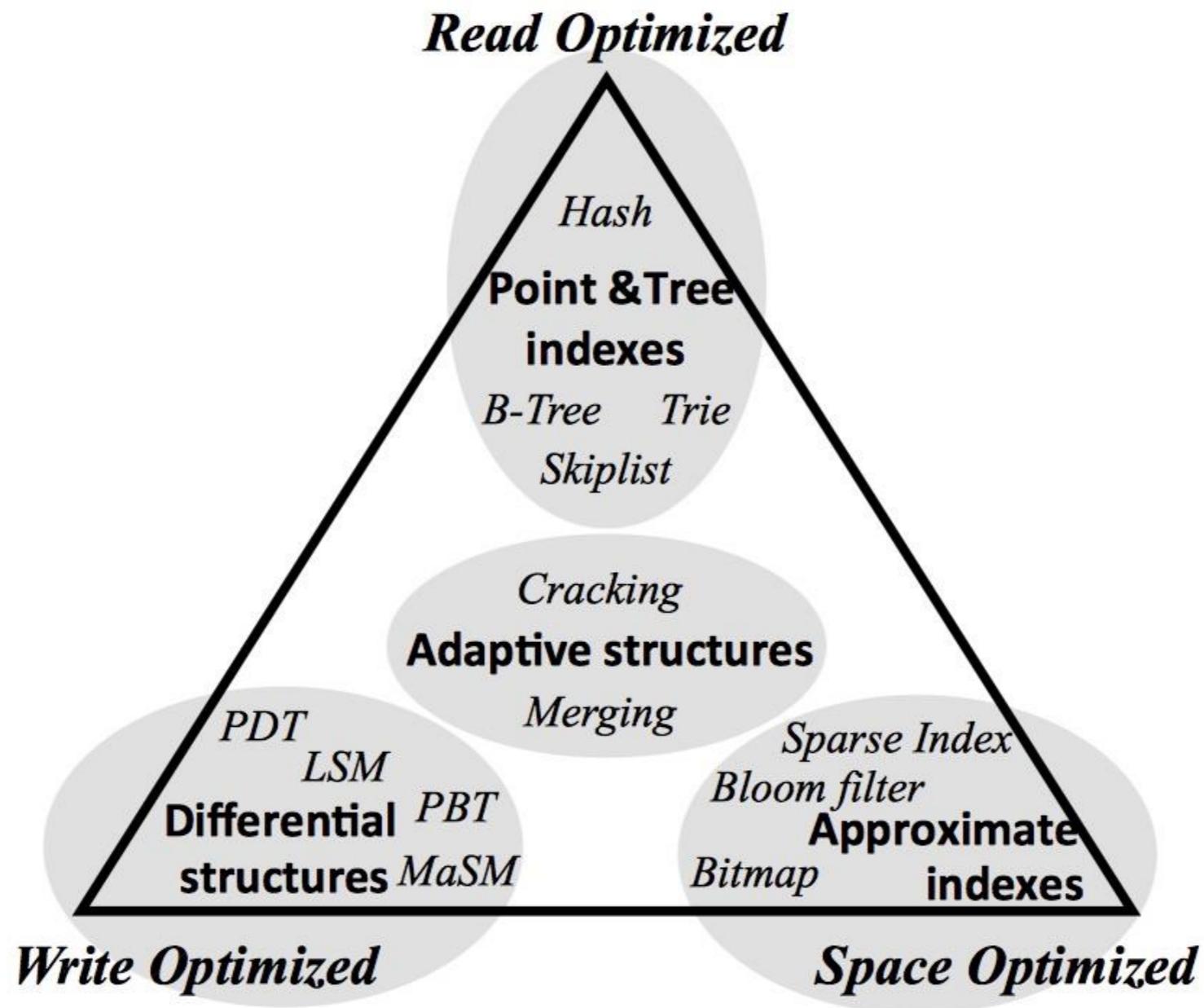


Figure 1: Popular data structures in the RUM space.

Data Structure	Write Amp (worst case)	Read Amp (range) (cold cache)	Read Amp (range) (warm)	Space Amp
B Tree	$O(B)$	$O(\log_B N/B)$	1	1.33
FT index	$O(k \log_k N/B)$	$O(\log_k N/B)$	1	negligible
LSM leveled	$O(k \log_k N/B)$	$O((\log^2 N/B)/\log k)$	3	2 (file-per-run) 1.1 (many files)
LSM size-tiered	$O(\log_k N/B)$	$O(k(\log^2 N/B)/\log k)$	13	3

O. 存储引擎

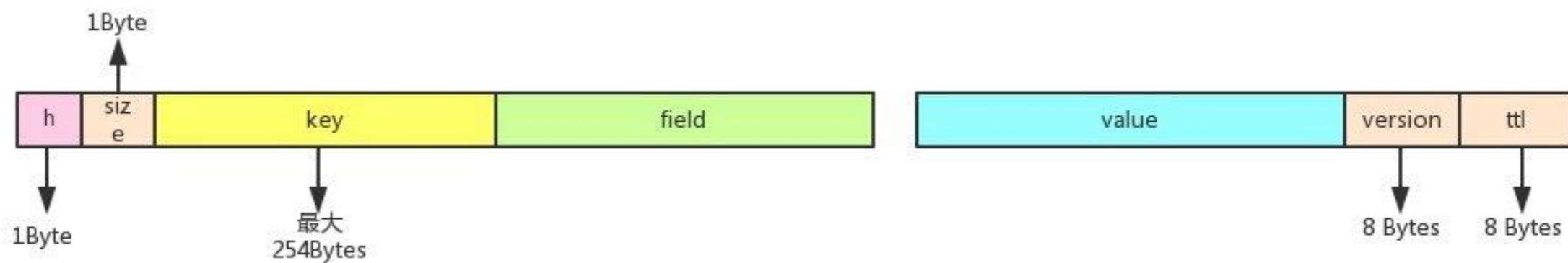


- Nemo
 - 基于rocksdb key-value 接口, 实现hash, list, set, zset 等接口
 - <https://github.com/Qihoo360/nemo>
- Meepo
 - Bitcask 存储引擎实现
 - <https://github.com/baotiao/meepo>

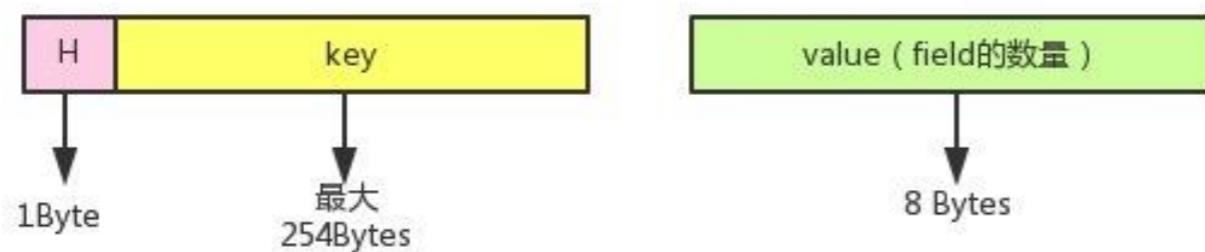
O. nemo-Hash



a. 每个hash键、field、value到落盘kv的映射转换



b. 每个hash键的元信息的落盘kv的存储格式



O.



zeppelin

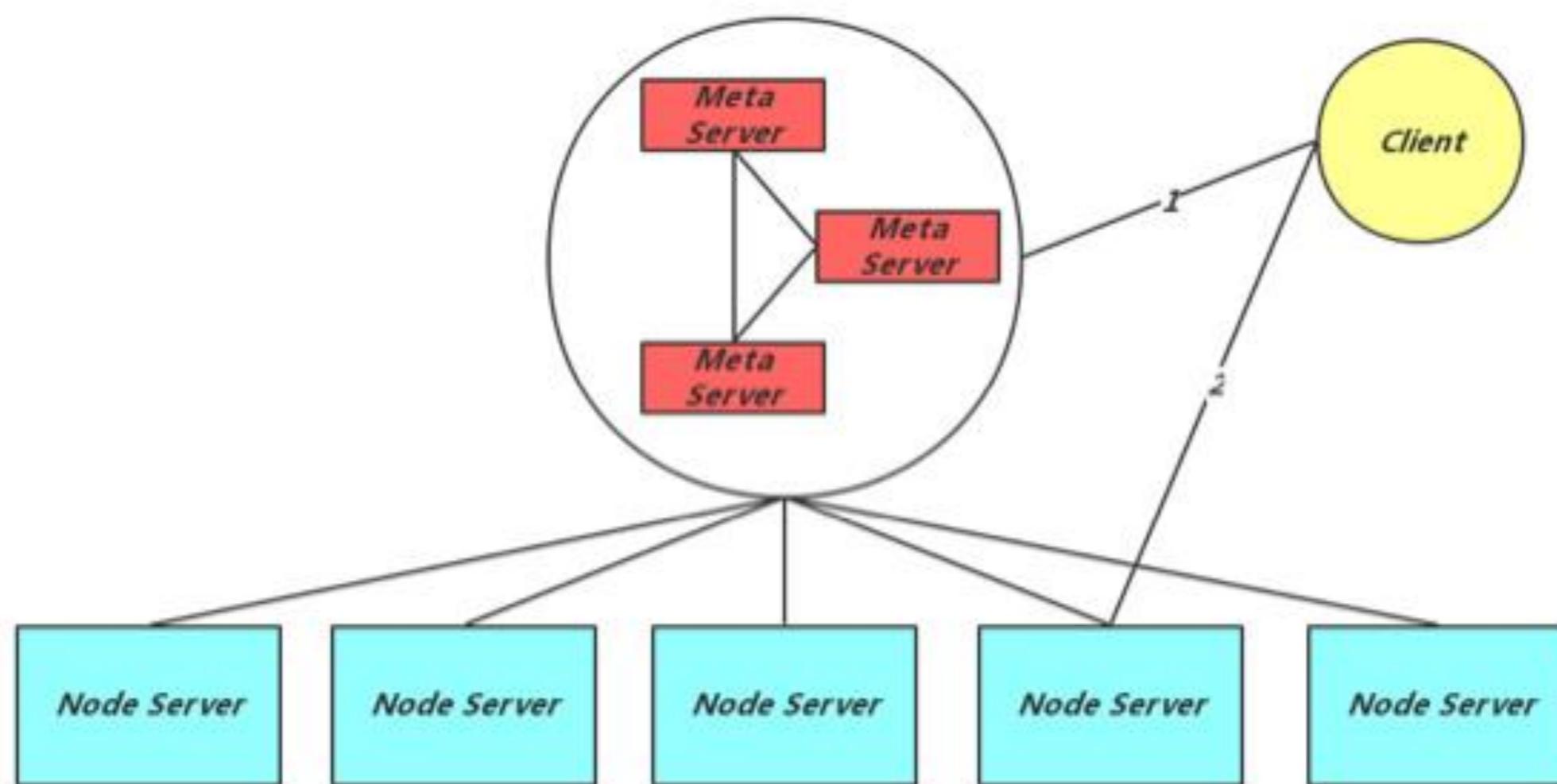
O. zeppelin



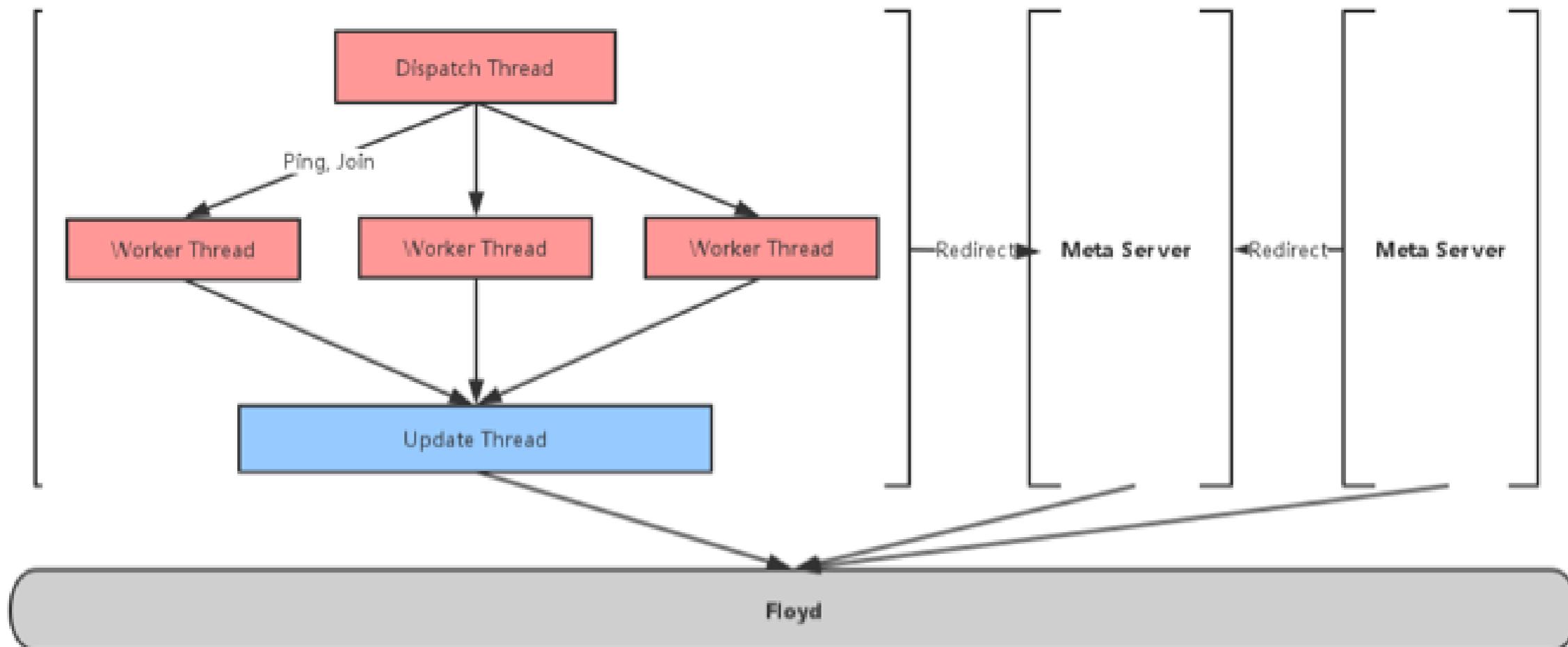
- 目标支持单机群千台
- 最终一致性
- 存储半离线数据

- 基础架构团队开发网络编程框架
- 高性能
 - 单节点支持80w
- 稳定
 - pika, zeppelin, gpstall, mongosync, 稳定使用3年
- 易用
 - <https://github.com/Qihoo360/pink>

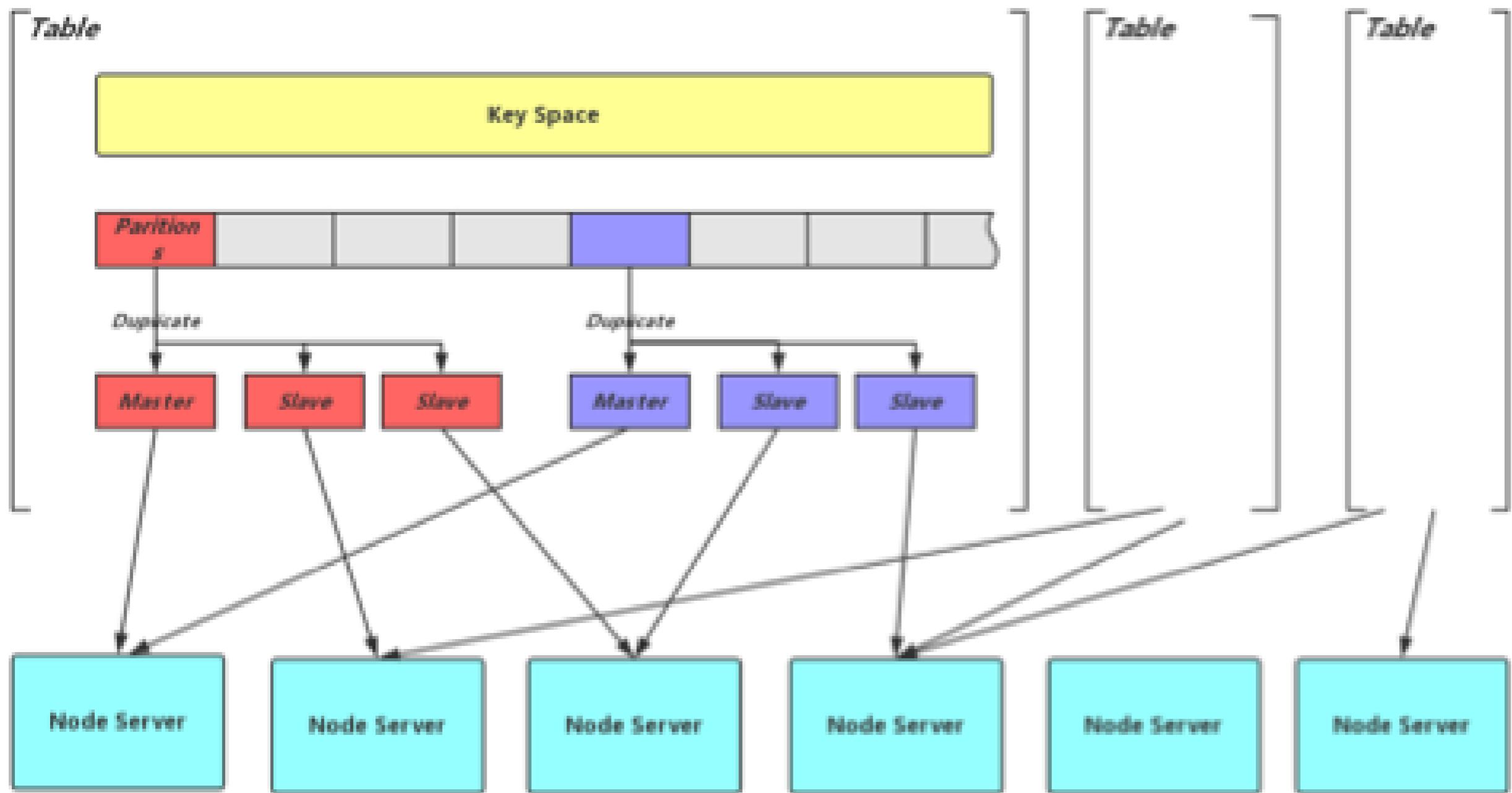
O. 整体结构

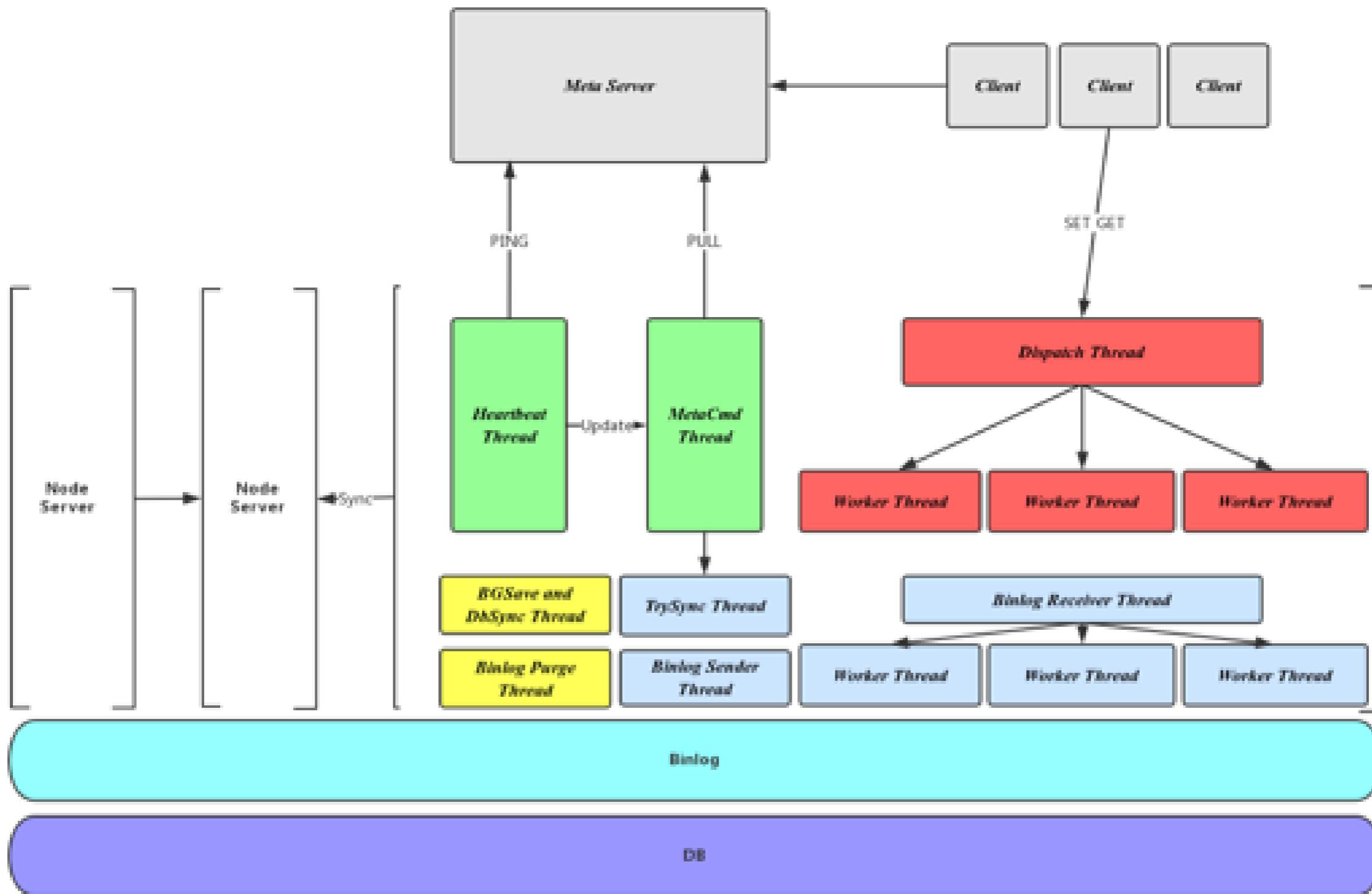


O. Meta 设计



O. 分表





O. 开源项目



- <https://github.com/Qihoo360>
- <https://github.com/Qihoo360/pika>
- <https://github.com/Qihoo360/pink>
- <https://github.com/baotiao/zeppelin>

O. Further reading:



- Ghemawat, Gobioff, & Leung. Google File System, SOSP 2003.
- G. DeCandia, et al., Dynamo: Amazon's Highly Available Key-Value Store
- B. Cooper, et al., PNUTS: Yahoo!'s Hosted Data Serving Platform
- L. Lamport, Paxos Made Simple
- Diego Ongaro, In Search of an Understandable Consensus Algorithm
- F. Chang, et al., Bigtable: A Distributed Storage System For Structured Data